

Shortest Unweighted Path

Remember that the (unweighted) length of a path in a graph is the number of edges the path contains. Consider a directed graph and let S be any specific node in this graph. We now give an algorithm for finding the shortest path from S to every other node in the graph.

The algorithm maintains a queue of nodes, which initially contains only S . It gives each node a value, which ultimately will be the length of the shortest path. Finally, it gives each node a predecessor node in its path from S .

Initially make the value of each node except S be "INFINITY". If you are thinking of a Java implementation, INFINITY can be either Integer.MAX_VALUE or Double.MAX_VALUE, depending on how you want to think of the values. Make the value of S be 0.

Now perform the following steps until the queue is empty.

- a) Remove the head of the queue. Call this node X .
- b) For each outgoing edge from X to another node Y , if the value of Y is INFINITY, make the new value of Y be the value of $X + 1$, make the predecessor of Y be X , and add Y to the queue.

Now, how do we know this algorithm works?

I claim that if node X has distance n from S then the value this algorithm assigns to X is n . This is certainly true when n is 0 or 1. For other nodes let $S=X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow \dots \rightarrow X_n=X$ be a path of length n to X . Suppose node X_t is the first node on this path that the algorithm assigns the wrong distance to. This means that node X_{t-1} has the correct distance. When X_{t-1} is removed from the queue and assigned the distance $t-1$, X_t will be added to the queue with distance t . So X_t in fact gets the correct distance. This means that all of the nodes in the path, including X , get the correct distance.

Note that this algorithm visits every edge in the graph (at least every edge that is reachable from S and so has running time $O(|V|)$).